ORIGINAL PAPER



The orbit of closure-involution operations: the case of Boolean functions

Jürgen Dassow¹

Received: 31 January 2020 / Accepted: 3 May 2021 / Published online: 15 May 2021 © The Author(s) 2021

Abstract

For a set *A* of Boolean functions, a closure operator *c* and an involution *i*, let $\mathcal{N}_{c,i}(A)$ be the number of sets which can be obtained from *A* by repeated applications of *c* and *i*. The orbit $\mathcal{O}(c, i)$ is defined as the set of all these numbers. We determine the orbits $\mathcal{O}(S, i)$ where *S* is the closure defined by superposition and *i* is the complement or the duality. For the negation non, the orbit $\mathcal{O}(S, \operatorname{non})$ is almost determined. Especially, we show that the orbit in all these cases contains at most seven numbers. Moreover, we present some closure operators where the orbit with respect to duality and negation is arbitrarily large.

Keywords Kuratowski's closure-complement theorem \cdot Superposition of Boolean functions \cdot Complement and negation and duality of sets of Boolean functions

Mathematics Subject Classification 06D25 · 06A15 · 08A05

1 Introduction

In Kuratowski (1922), proved the following closure-complement theorem: If (X, T) is a topological space and $A \subseteq X$, then at most 14 sets can be obtained from A by repeated applications of the operations topological closure and complement. Furthermore, there is a topological space and a set for which the bound 14 is achieved. More information on Kuratowski-like theorems for topological space can be found in Gardner and Jackson (2008).

Hammer (1960) noticed that such a statement holds in a more general setting; it is not necessary to consider topological spaces and topological closure. The theorem

Dedicated to Prof. Gustav Burosch on the Occasion of his 80th Birthday.

☑ Jürgen Dassow dassow@iws.cs.uni-magedeburg.de

¹ Fakultät für Informatik, Otto-von-Guericke-Universität Magdeburg, PSF 4120, 39106 Magdeburg, Germany

also holds if—instead of the topological closure—a closure operator on a set X is used.

Most papers related to Kuratowski's closure-complement theorem ask for upper bounds for the number of sets obtainable by repeated applications of a closure operator and complement. However, one can also consider the following more general question: Given a closure operator, determine the set of all numbers n (called the orbit of the closure operator and complement) such that there is a set A_n from which we can obtain exactly n sets by repeated applications of the closure operator and complement.

In Brzozowski et al. (2009), this question was firstly investigated by Brzozowski, Grant, and Shallit for the Kleene-closure defined on formal languages and complement. They proved that the orbit of Kleene-closure and complement consists of the numbers 4, 6, 8, 10, 12, 14. Moreover, for $n \in \{4, 6, 8, 10, 12, 14\}$, they gave precise conditions for a language to produce exactly *n* languages by repeated applications of Kleene-closure and complement.

For further language theoretic closure operators and involutions (instead of complement) the orbit was studied in Dassow (2019).

In this paper, we continue the determination of the orbit, but we consider the set of Boolean functions. Here a "classical" closure operator is defined by superpositions. The lattice of all closed sets (with respect to superpositions) of Boolean functions was determined in 1921 by Post (see Post 1921, a more complete version is Post 1941, and a modern version is Jablonski et al. 1970). We study the orbit of superpositions and complement, duality, and negation as involution. We prove that the orbit contains three, four, and at most 7 numbers for complement, duality, and negation, respectively. The corresponding Kuratowski numbers are six, four, and seven, respectively.

However, the situation changes completely if we allow other closure operators. We define some special closure operators such that with duality (or negation) the corresponding orbit contains infinity or has *m* elements where *m* is an arbitrary natural number with $m \ge 3$.

2 Definitions and known facts

Let *X* be a set. We define the complement \overline{A} of a set $A \subset X$ by $\overline{A} = X \setminus A$.

An operator c is called a *closure operator* on X, if the following three conditions are satisfied:

- For all sets $A \subseteq X$, $c(A) \subseteq X$.
- For all sets $A \subseteq X$, $A \subseteq c(A)$.
- For all sets $A \subseteq X$ and $B \subseteq X$, $A \subseteq B$ implies $c(A) \subseteq c(B)$.
- For all sets $A \subseteq X$, c(c(A)) = c(A), i.e., c is idempotent.

An operator *i* is called an *involution* on *X* if, for any $A \subseteq X$, the relations $i(A) \subseteq X$ and i(i(A)) = A hold.

Definition 1 Let *c* a closure operator on *X*, and *i* an involution on *X*. Then, for $A \subseteq X$, we define the orbit $\mathcal{O}_{c_i}^X(A)$ of *A* under *c* and *i* as the set of all sets which can be obtained

from A by repeated applications of c and i and set

$$\mathcal{N}_{c,i}^X(A) = \operatorname{card}(\mathcal{O}_{c,i}(A)).$$

Moreover, we define the orbit of the pair (c, i) as

$$\mathcal{O}^X(c,i) = \{n \mid \mathcal{N}_{c,i}^X(A) = n \text{ for some } A \subseteq X\}$$

and the Kuratowski number of (c, i) as

$$\mathcal{K}^X(c,i) = \sup\{n \mid n \in \mathcal{O}^X(c,i)\}.$$

In this terminology, the classical Kuratowski's Theorem is given as follows: If (X, \mathcal{T}) is a topological space and c the corresponding topological closure, then $\mathcal{K}^X(c, -) = 14$.

By the properties of an closure operator *c* and an involution *i*, in order to determine the orbit $\mathcal{O}_{c,i}^X(A)$ of a set *A* it is sufficient to determine the sets

$$A, c(A), i(c(A)), c(i(c(A))), i(c(i(c(A)))), \dots$$

and

$$A, i(A), c(i(A)), i(c(i(A))), c(i(c(i(A)))), \dots$$

Let P_2 be the set of Boolean functions. In the rest of the paper we restrict to $X = P_2$, and for the sake of simplicity we shall omit the upper index P_2 in the notations given in Definition 1.

We now define some special Boolean functions and sets of Boolean functions, which will be used later. With respect to the notation, we follow (Jablonski et al. 1970).

The constants giving the value 0 or 1 are denoted by k_0 and k_1 , respectively.

A special unary Boolean function is the negation non defined by

non(0) = 1 and non(1) = 0. We extend the negation to functions by setting $(non(f))(x_1, x_2, ..., x_n) = non(f(x_1, x_2, ..., x_n)).$

We use the following functions (in some cases we give two notations, and if the functions are associative, we omit some brackets in the sequel):

• $\operatorname{vel}(x_1, x_2) = x_1 \lor x_2 = 0$ if and only if $x_1 = x_2 = 0$,

• $g^k(x_1, x_2, ..., x_k) = x_1 \lor x_2 \lor \cdots \lor x_k$ for $k \ge 2$,

• $et(x_1, x_2) = x_1 \cdot x_2 = 1$ if and only if $x_1 = x_2 = 1$,

- $h^k(x_1, x_2, ..., x_k) = x_1 \cdot x_2 \cdot \cdots \cdot x_k$ for $k \ge 2$,
- $x_1 + x_2 = 0$ if and only if $x_1 = x_2$,

• $\operatorname{sh}(x_1, x_2) = \operatorname{non}(\operatorname{vel}(x_1, x_2))$ and $\operatorname{sh}'(x_1, x_2) = \operatorname{non}(\operatorname{et}(x_1, x_2))$.

Let C_2 (C_3) be the sets of all functions f such that f(1, 1, ..., 1) = 1(f(0, 0, ..., 0) = 0, respectively).

The dual function d(f) of a function f is defined as

 $(d(f))(x_1, x_2, \dots, x_n) = (\operatorname{non}(f))(\operatorname{non}(x_1), \operatorname{non}(x_2), \dots, \operatorname{non}(x_n)).$

Moreover, a Boolean function f is called self-dual if and only if d(f) = f. Let D_3 denote the set of all self-dual Boolean functions.

We extend the concept of negation and duality to subsets $A \subseteq P_2$ by setting

$$\operatorname{non}(A) = \{\operatorname{non}(f) \mid f \in A\} \text{ and } d(A) = \{d(f) \mid f \in A\}.$$

Note that these operators non and d are involutions on P_2 .

A function *f* is called linear if $f(x_1, x_2, ..., x_n) = x_{i_1} + x_{i_2} + \cdots + x_{i_r}$ or $f(x_1, x_2, ..., x_n) = x_{i_1} + x_{i_2} + \cdots + x_{i_r} + 1$, where $1 \le i_1 < i_2 < \cdots < i_r \le n$. By L_1 we denote the family of all linear functions.

By 0 < 1, an order is defined on $\{0, 1\}$. We say that an *n*-ary function *f* is monotone if and only if, for all tuples $(x_1, x_2, ..., x_n)$ and $(y_1, y_2, ..., y_n), x_i \le y_i$ for $1 \le i \le n$ implies $f(x_1, x_2, ..., x_n) \le f(y_1, y_2, ..., y_n)$.

We now define some operations which lead to the closure operator superposition. For an *n*-ary function f, $n \ge 0$, we set

$$(\zeta_0(f))(x_{n+1}, x_1, x_2, \dots, x_n) = f(x_1, x_2, \dots, x_n)$$
(1)

$$(\zeta_i(f))(x_1, x_2, \dots, x_i, x_{n+1}, x_{i+1}, \dots, x_n) = f(x_1, x_2, \dots, x_n), \ 1 \le i \le n.$$
(2)

For an *n*-ary function $f, n \ge 2$, and a permutation π on $\{1, 2, ..., n\}$, we define

$$(\Delta(f))(x_1, x_2, \dots, x_{n-1}) = f(x_1, x_2, \dots, x_{n-1}, x_{n-1}),$$
(3)

$$(\pi(f))(x_1, x_2, \dots, x_n) = f(x_{\pi(1)}, x_{\pi(2)}, \dots, x_{\pi(n)}).$$
(4)

If f is an n-ary function and $n \le 1$, then we set $\Delta(f) = \pi(f) = f$.

For an *n*-ary function $f, n \ge 1$, and an *m*-ary function $g, m \ge 0$, we define the (n + m - 1)-ary function

$$(f \circ g)(x_1, \dots, x_{n-1}, y_1, \dots, y_m) = f(x_1, \dots, x_{n-1}, g(y_1, \dots, y_m)).$$
 (5)

For a set $A \subseteq P_2$, we define [A] as the set of all functions which can be obtained by finitely many iterated application of the operations defined in (1) – (5) to functions from A. It is easy to see that the operator S given by S(A) = [A] is a closure operator. Thus, A is called closed if and only if A = [A].

We denote the set of all functions which can be obtained from f by iterated application of (1) and (2) by $\ll f \gg$.

3 The classical closure operator: superpositions

In this section we study the orbits of the closure operator given by superpositions and the involutions complement, duality, and negation.

Theorem 1 We have $\mathcal{O}(S,^{-}) = \{2, 4, 6\}.$

Proof Let A be a subset of P_2 . First, we assume that $[\underline{A}] \neq P_2$. Then sh $\notin A$ and sh $\notin [A]$ (since $[\{sh\}] = P_2$). Therefore sh $\in \overline{A}$, sh $\in \overline{[A]}$, and $[\overline{A}] = [\overline{[A]}] = P_2$. Moreover $\overline{P_2} = \emptyset$, $[\emptyset] = \emptyset$, and $[P_2] = P_2$. Consequently, $\mathcal{O}_{S,-}(A)$ contains at most the six sets $A, \overline{A}, [A], \overline{[A]}, P_2$, and \emptyset .

Now let $[A] = P_2$. If $[\overline{A}] = P_2$, too, then $\mathcal{O}_{S,-}(A)$ contains at most the sets A, \overline{A} , P_2 , and \emptyset . If $[\overline{A}] \neq P_2$, then sh $\in [\overline{A}]$, which results in $[[\overline{A}]] = P_2$; and consequently at most the six sets $A, \overline{A}, [\overline{A}], [\overline{A}], P_2$, and \emptyset are in $\mathcal{O}_{S,-}(A)$.

Furthermore, if $\mathcal{O}_{S,-}(A)$ contains a set *B*, then it also contains \overline{B} . Hence $\mathcal{O}(S,-)$ contains only even numbers.

Consequently, the only possible numbers which can occur in $\mathcal{O}(S, -)$ are 2,4, and 6. We now prove that all these numbers are possible.

Let $B_1 = \{\text{vel}\}$. Then $[B_1] = \bigcup_{k \ge 1} \ll g^k \gg \text{ and thus (as shown above) } [\overline{B_1}] = [\overline{[B_1]}] = P_2$. Consequently,

$$\mathcal{O}_{S,-}(B_1) = \{B_1, \overline{B_1}, [B_1], \overline{[B_1]}, P_2, \emptyset\}.$$

Obviously, the six sets of $\mathcal{O}_{S,-}(B_1)$ are pairwise different. Hence $6 \in \mathcal{O}(S,-)$.

Let $B_2 = \{\text{sh}\}$. Then $\text{sh}' \in \overline{B_2}$ and $[B_2] = [\overline{B_2}] = P_2$. Thus

$$\mathcal{O}_{S,-}(B_2) = \{B_2, B_2, P_2, \emptyset\}.$$

Since these sets are pairwise different, $4 \in \mathcal{O}(S, -)$.

Let $B_3 = P_2$. Then $\mathcal{O}_{S,-}(B_3) = \{P_2, \emptyset\}$, and consequently $2 \in \mathcal{O}(S,-)$.

Theorem 2 We have $O(S, d) = \{1, 2, 3, 4\}$.

Proof Let A be a subset of P_2 . Because d([A]) = [d(A)] (Jablonski et al. 1970, Chapter 1, Sect. 6, Folgerung 2), we obtain that $\mathcal{O}_{S,d}(A)$ contains at most the sets A, d(A), [A], and [d(A)]. Thus the only possible numbers in $\mathcal{O}(S, d)$ are 1, 2, 3, and 4. We now give witnesses for these numbers.

Let $B_1 = P_2$. Then $\mathcal{O}_{S,d}(B_1) = \{P_2\}$ and, consequently, $1 \in \mathcal{O}(S, d)$.

Let $B_2 = \{g\}$ with $g(x, y, z) = x \cdot \operatorname{non}(y) \lor x \cdot \operatorname{non}(z) \lor \operatorname{non}(y) \lor z$. Then it is known (see Jablonski et al. 1970) that g is a self-dual function and $[\{g\}]$ is the set D_3 of all self-dual functions. Thus we obtain $d(\{g\}) = \{g\}$ and $d([\{g\}]) = [\{g\}]$. This implies $\mathcal{O}_{S,d}(B_2) = \{B_2, [B_2]\}$ and $2 \in \mathcal{O}(S, d)$.

Let $B_3 = \{\text{vel, non}\}$. Then we have $d(B_3) = \{\text{et, non}\}$ and $[\{\text{vel, non}\}] = [\{\text{et, non}\}] = P_2$ which gives $\mathcal{O}_{S,d}(B_3) = \{B_3, d(B_3), P_2\}$. Thus $3 \in \mathcal{O}(S, d)$.

Let $B_4 = \{\text{vel}\}$. Then $d(B_4) = \{\text{et}\}, [B_4] = \bigcup_{k \ge 1} \ll g^k \gg$, and $[d(B_4)] = \bigcup_{k \ge 1} \ll h^k \gg$. Therefore $4 \in \mathcal{O}(S, d)$.

Theorem 3 We have $\{1, 2, 3, 4, 5, 7\} \subseteq \mathcal{O}(S, \text{non}) \subseteq \{1, 2, 3, 4, 5, 6, 7\}$.

Proof Let A be a subset of P_2 .

Assume that $[A] = P_2$. If $[non(A)] = P_2$ holds, then we have $\mathcal{O}_{S,non}(A) = \{A, non(A), P_2\}$ (we do not know whether these sets are pairwise different or some of them are identical). If $[non(A)] \neq P_2$, then $A \subseteq non([non(A)])$ and hence

 $[\operatorname{non}([\operatorname{non}(A)])] = P_2$. Thus $\mathcal{O}_{S,\operatorname{non}}(A)$ contains at most the sets A, $\operatorname{non}(A)$, $[\operatorname{non}(A)]$, $\operatorname{non}([\operatorname{non}(A)])$, and P_2 .

Now we discuss some cases where $[A] \neq P_2$. We start with the cases that $[A] = O_i$ for some *i* with $1 \le i \le 9$ where

$$O_{1} = \ll \mathrm{id} \gg \qquad O_{5} = \ll \mathrm{id} \gg \cup \ll k_{1} \gg$$

$$O_{2} = \ll k_{1} \gg \qquad O_{6} = \ll \mathrm{id} \gg \cup \ll k_{0} \gg$$

$$O_{3} = \ll k_{0} \gg \qquad O_{7} = \ll k_{1} \gg \cup \ll k_{0} \gg$$

$$O_{4} = \ll \mathrm{id} \gg \cup \ll \mathrm{non} \gg \qquad O_{8} = \ll \mathrm{id} \gg \cup \ll k_{1} \gg \cup \ll k_{0} \gg$$

$$O_{9} = \ll \mathrm{id} \gg \cup \ll \mathrm{non} \gg \cup \ll k_{1} \gg \cup \ll k_{0} \gg$$

If $[A] = O_1$, then $A \subseteq O_1$ and non $(A) \subseteq \ll$ non \gg . Therefore $[non(A)] = O_4$. Furthermore, non $([A]) = non(O_1) = \ll$ non \gg and $[non([A])] = O_4$. Because $[O_4] = non(O_4) = O_4$ we have that $\mathcal{O}_{S,non}(A)$ contains at most the sets A, non(A), O_1 , \ll non \gg , and O_4 .

If $[A] = O_2$, then $A \subseteq O_2$, non $(A) \subseteq O_3$, and $[non(A)] = O_3$. Because non $(O_2) = O_3$, non $(O_3) = O_2$, $[O_2] = O_2$, and $[O_3] = O_3$, we have that $\mathcal{O}_{S,non}(A)$ contains at most the sets A, non(A), O_2 , and O_3 .

If $[A] = O_3$, we get analogously that $\mathcal{O}_{S,non}(A)$ contains at most the sets A, non(A), O_2 , and O_3 .

If $[A] = O_4$, then A contains at least one function f_1 of \ll non \gg . Assume that $A \subseteq \ll$ non \gg . Then we get the relations non $(A) \subseteq O_1$, $[non(A)] = O_1$, non $([non(A)]) = \ll$ non \gg , and $[non([non(A)])] = O_4$. Because non $(O_4) = [O_4] = O_4$, $\mathcal{O}_{S,non}(A)$ contains at most the sets A, non(A), O_1 , \ll non \gg , and O_4 . If we assume that A contains a function of O_1 , too, then non(A) contains a function of \ll non \gg and consequently $[non(A)] = O_4$. Therefore $\mathcal{O}_{S,non}(A)$ contains at most the sets A, non(A), and O_4 .

If $[A] = O_5$, then A contains functions $f_1 \in \ll id \gg and f_2 \in \ll k_1 \gg$. Then non(A) contains a function of $\ll non \gg and a$ function of the set $\ll k_0 \gg$. This implies $[non(A)] = O_9$. Because we have that $non(O_5) = \ll non \gg \cup \ll k_0 \gg$, we also obtain $[non(O_5)] = O_9$. Therefore, $\mathcal{O}_{S,non}(A)$ contains at most the sets A, non(A), O_5 , non(O₅) and O₉.

Analogously, if $[A] = O_6$, then $\mathcal{O}_{S,\text{non}}(A)$ contains at most the sets A, non(A), O_6 , non(O_6) and O_9 .

If $[A] = O_7$, A contains a function of $\ll k_0 \gg$ and a function of $\ll k_1 \gg$. This statement holds for non(A), too. Therefore $[non(A)] = O_7$. Because $[O_7] = non(O_7) = O_7$, $\mathcal{O}_{S,non}(A)$ contains at most the sets A, non(A), and O_7 .

If $[A] = O_8$, A contains functions $f_1 \in \ll \operatorname{id} \gg$ and $f_2 \in \ll k_1 \gg$, and $f_3 \in \ll k_0 \gg$. Therefore non(A) and non([A]) contain non $(f_1) \in \ll$ non \gg and non $(f_2) \in \ll k_0 \gg$. Hence $[\operatorname{non}(A)] = [\operatorname{non}([A])] = O_9$. Thus $\mathcal{O}_{S,\operatorname{non}}(A)$ contains at most the sets A, non(A), O_8 , and O_9 .

If $[A] = O_9$, the set A contains a function $f_1 \in \ll \operatorname{non} \gg$ and a function $f_2 \in \ll k_i \gg$ for some $i \in \{0, 1\}$. If A contains $f_3 \in \ll \operatorname{id} \gg$, then we

get $[non(A)] = O_9$. Otherwise, $non(A) \subseteq O_8$ and $[non(A)] \in \{O_5, O_6, O_8\}$. In all these three cases non([non(A)]) contains a function of $\ll non \gg$ and a function of

 $\ll k_i \gg$ for some $i \in \{0, 1\}$. Consequently, [non([non(A)])] equals O_9 . Therefore, $\mathcal{O}_{S,non}(A)$ contains at most the sets A, non(A), [non(A)], non([non(A)]), and O_9 .

We now discuss the case that $\{\text{vel}, \text{et}\} \cap [A] \neq \emptyset$. Then we obtain the relation $\{\text{sh}, \text{sh'}\} \cap \text{non}([A]) \neq \emptyset$ which implies $[\text{non}([A])] = P_2$. Moreover, A contains a non-self-dual function f_1 and a non-linear function f_2 (otherwise [A] would contain only self-dual or only linear functions, which cannot hold by our assumption because vel and et are neither self-dual nor linear. But then $\text{non}(f_1)$ is not self-dual and $\text{non}(f_2)$ is not linear. Thus [A] is not contained in D_3 and not contained in L_1 . Moreover, A has to contain a function f_3 which depends essentially on at least two variables (otherwise only functions in \ll id $\gg \cup \ll$ non $\gg \cup \ll k_0 \gg \cup \ll k_1 \gg$ could be generated, which gives $\{\text{vel}, \text{et}\} \cap [A] = \emptyset$ in contrast to our assumption). Obviously, $\text{non}(f_3)$ also depends on at least two variables. Thus [non(A)] is not contained in O_9 .

By the Post graph of all closed sets (see Jablonski et al. 1970, page 76), [non(A)] has to contain a function of $\bigcup_{k\geq 1} \ll g^k \gg$ or a function of $\bigcup_{k\geq 1} \ll h^k \gg$. From this it follows that $\{vel, et\} \cap [non(A)]$ is not empty. As above, we get that [non([non(A)])] = P_2 . Hence $\mathcal{O}_{S,non}(A)$ contains at most the sets A, non(A), [A], non([A]), [non(A)], non([non(A)]), and P_2 .

We now turn to $[A] \in \{L_1, L_2, L_3, L_4, L_5\}$ where

$$L_2 = L_1 \cap C_2, \ L_3 = L_1 \cap C_3, \ L_4 = L_1 \cap C_2 \cap C_3, \ L_5 = L_1 \cap D_3.$$

We note that all functions in L_4 are self-dual.

Let $[A] = L_1$. Then we have that $A \subseteq \text{non}[\text{non}(A)]$. Consequently, we obtain $[\text{non}[\text{non}(A)] = L_1$. Therefore $\mathcal{O}_{S,\text{non}}(A)$ contains at most the sets A, non(A), [non(A)], non([non(A)]), and L_1 .

Let $[A] = L_2$. Then the set A contains a non-self-dual function g. Moreover, id $\in L_2$. Hence non (L_2) contains non $(g) \notin D_3$ and non $\notin C_2 \cup C_3$. Thus we obtain $[non(L_2)] = L_1$. Furthermore, $A \subseteq non([non(A)])$ which implies the relation $[non([non(A)])] \in \{L_2, L_1\}$. From the above considerations, we obtain that $\mathcal{O}_{S,non}(A)$ contains at most the sets A, non(A), [non(A)], non([non(A)]), L_2 , non (L_2) , and L_1 .

The case $[A] = L_3$ can be handled analogously, and gives also at most seven sets in $\mathcal{O}_{S,\text{non}}(A)$.

Let $[A] = L_4$. Since id $\in L_4$, we get non $\in \text{non}(L_4)$ and $[\text{non}(L_4)] = L_5$. Moreover, $A \subseteq \text{non}([\text{non}(A)])$ which implies $[\text{non}([\text{non}(A)])] \in \{L_4, L_5\}$ and that $\mathcal{O}_{S,\text{non}}(A)$ contains at most the sets A, non(A), [non(A)], non([non(A)]), L_4 , non (L_4) , and L_5 .

Let $[A] = L_5$. Because $A \subseteq \text{non}([\text{non}(A)])$, we obtain $[\text{non}([\text{non}(A)])] = L_5$. Thus $\mathcal{O}_{S,\text{non}}(A)$ contains at most the sets A, non(A), [non(A)], non([non(A)]), and L_5 .

Finally, we discuss the case $[A] \in \{D_1, D_2, D_3\}$, where D_1 is the set of all monotone self-dual functions and $D_2 = D_3 \cap C_2 \cap C_3$.

Let $[A] = D_3$. Because $A \subseteq \text{non}[\text{non}(A)]$, we obtain $[\text{non}[\text{non}(A)] = D_3$ and that $\mathcal{O}_{S,\text{non}}(A)$ contains at most the sets A, non(A), [non(A)], non([non(A)]), and D_3 .

Let $[A] = D_1$ or $[A] = D_2$. Then id $\in [A]$, and therefore non in non([A]), but it is not monotone, not in C_2 and not in C_3 . Moreover, A and thus non(A) contain

a non-linear function each. Moreover, for any function $f \in C_2 \cap C_3$, we have that $\operatorname{non}(f) \notin C_2 \cup C_3$ and $\operatorname{non}(f)$ is not monotone. Since $D_2 \subset D_1 \subset C_2 \cap C_3$, again, by the Post graph of all closed sets, we get $[\operatorname{non}(A)] = D_3$. Moreover, $[\operatorname{non}([A])] = D_3$ now follows immediately. Hence $\mathcal{O}_{S,\operatorname{non}}(A)$ contains at most the sets A, $\operatorname{non}(A)$, [A], $\operatorname{non}([A])$, and D_3 .

By Jablonski et al. (1970), we have above covered all possible cases.

Summarizing, in all cases $\mathcal{O}_{S,non}(A)$ contains at most seven sets. We now show that the numbers 1,2,3,4,5, and 7 are possible as the cardinality of some orbit $\mathcal{O}_{S,non}(A)$.

For P_2 , obviously, we get $\mathcal{O}_{S,\text{non}}(P_2) = \{P_2\}$. Thus $1 \in \mathcal{O}(S, \text{non})$.

For $A_2 = \{$ sh, vel $\}$, we obtain non $(A_2) = A_2$ and $[A_2] = P_2$. Consequently, $\mathcal{O}_{S,\text{non}}(A_2) = \{A_2, P_2\}$ and hence $2 \in \mathcal{O}(S, \text{non})$.

For $A_3 = \{\text{vel, non}\}$, we have $\text{non}(A_3) = \{\text{sh, id}\}$. Furthermore, we get $[A_3] = [\text{non}(A_3)] = P_2$, which implies $\mathcal{O}_{S,\text{non}}(A_3) = \{A_3, \text{non}(A_3), P_2\}$ and $3 \in \mathcal{O}(S, \text{non})$.

For $A_4 = \{k_0\}$, we obtain the relations $non(A_4) = \{k_1\}, [A_4] = \ll k_0 \gg$, and $[non(A_4)] = \ll k_1 \gg$. Therefore

$$\mathcal{O}_{S,\text{non}}(A_4) = \{\{k_0\}, \{k_1\}, \ll k_0 \gg, \ll k_1 \gg\}$$

and $4 \in \mathcal{O}(S, \text{non})$ hold.

For $A_5 = \{id, k_1\}$, we get the relations $\operatorname{non}(A_5) = \{\operatorname{non}, k_0\}$, $[A_5] = O_5$, $\operatorname{non}(O_5) = \ll \operatorname{non} \gg \cup \ll k_0 \gg$, and $[\operatorname{non}(A_5)] = [\operatorname{non}(O_5)] = O_9$. Therefore we have

 $\mathcal{O}_{S,\text{non}}(A_4) = \{\{\text{id}, k_1\}, \{\text{non}, k_0\}, O_5, \ll \text{non} \gg \cup \ll k_0 \gg, O_9\}$

and $5 \in \mathcal{O}(S, \text{non})$.

Let $A_7 = \{g\}$ where $g(x, y) = x \vee \operatorname{non}(y)$. We recall that F_4^{∞} (and F_8^{∞}) is the set of all functions f satisfying that there is an i such that $f^{-1}(0) \subseteq \{0, 1\}^{i-1} \times \{0\} \times \{0, 1\}^{n-i}$ (and $f^{-1}(1) \subseteq \{0, 1\}^{i-1} \times \{1\} \times \{0, 1\}^{n-i}$, respectively). By Jablonski et al. (1970), $[\{g\}] = F_4^{\infty}$. Let $g' = \operatorname{non}(g)$. Then $g'(x, y) = \operatorname{non}(x) \wedge y$ and $[\{g'\}] = F_8^{\infty}$ (because the function $g''(x, y) = x \wedge \operatorname{non}(y)$, which is a generator of F_8^{∞} , is obtained from g' by a permutation of variables and vice versa). Obviously, the function h with $h(x, y, z) = x \wedge (y \vee z)$ is in F_8^{∞} , but $\operatorname{non}(h)$ is neither in F_8^{∞} nor in F_4^{∞} . By definition, for any n-ary function from the set $\operatorname{non}(F_4^{\infty})$, there is a j such that $f^{-1}(1) \subseteq \{0, 1\}^{j-1} \times \{0\} \times \{0, 1\}^{n-j}$. Now it is obvious that $\operatorname{non}(h)$ is not in $\operatorname{non}(F_4^{\infty})$. By these facts and the infinity of F_8^{∞} , $\operatorname{non}(F_8^{\infty})$ differs from A_7 , $\operatorname{non}(A_7)$, F_4^{∞} , F_8^{∞} , and $\operatorname{non}(F_8^{\infty})$. Moreover, vel $\in F_4^{\infty}$ and et $\in F_8^{\infty}$. Therefore sh $\in \operatorname{non}(F_4^{\infty})$, sh' $\in \operatorname{non}(F_8^{\infty})$, and $[\operatorname{non}(F_4^{\infty})] = [\operatorname{non}(F_8^{\infty})] = P_2$. Therefore we get

$$\mathcal{O}_{S,\text{non}}(A_7) = \{A_7, \text{non}(A_7), F_4^{\infty}, \text{non}(F_4^{\infty}), F_8^{\infty}, \text{non}(F_8^{\infty}), P_2\}$$

and $7 \in \mathcal{O}(S, \operatorname{non})$.

It remains as an open problem whether $6 \in \mathcal{O}(S, \text{ non})$. We conjecture that six does not belong to $\mathcal{O}(S, \text{ non})$. The reason for that is that there are only a few cases where

six can occur (mostly we got that at most five sets are in the orbit of A), and for some of them we can show that six is impossible.

Corollary 1 With respect to the Kuratowski number the following relations hold:

$$\mathcal{K}(S,^{-}) = 6, \ \mathcal{K}(S,d) = 4, \ and \ \mathcal{K}(S,\operatorname{non}) = 7.$$

4 Special closure operators

In the preceding section, we have studied the orbit of superposition and some involutions. In all cases, the Kuratowski number is smaller than 7 and therefore we only get very small orbits. We shall now prove that this depends on the closure operator superposition. If we consider other closure operators on sets of Boolean functions and the involutions duality or negation, we can obtain arbitrary large Kuratowski numbers and arbitrary large orbits.

Theorem 4 *There is a closure operation* c_1 *such that* $\mathcal{O}(c_1, d) = \{1, 2, 3, \infty\}$ *.*

Proof We say that $f \in P_2$ is a β -function (or γ -function), if f(x, ..., x) = 1(f(x, ..., x) = 0, respectively) for $x \in \{0, 1\}$. For $n \ge 0$ and $z \in \{\beta, \gamma\}$, by U(z, n), we denote the set of all *n*-ary *z*-functions. Obviously, $U(\beta, 0) = \{k_1\}$ and $U(\gamma, 0) = \{k_0\}$. Hence, for $\kappa \in \{\beta, \gamma\}$, there is no *A* such that $\emptyset \subset A \subset U(\kappa, 0)$. Moreover, we set

$$V(\beta, n) = \bigcup_{i=0}^{n} U(\beta, i) \text{ and } V(\gamma, n) = \bigcup_{i=0}^{n} U(\gamma, i).$$

We note that the dual of a β -function is a γ -function, and vice versa. Thus $d(U(\beta, n)) = U(\gamma, n)$ and $d(U(\gamma, n)) = U(\beta, n)$ for $n \ge 0$.

We define c_1 by

$$c_{1}(A) = \begin{cases} \emptyset & \text{for } A = \emptyset \\ V(\beta, n+1) & \text{for } A \subseteq V(\beta, n), \ A \cap U(\beta, n) \neq \emptyset, \ n \text{ even} \\ V(\beta, n) & \text{for } A \subseteq V(\beta, n), \ A \cap U(\beta, n) \neq \emptyset, \ n \text{ odd} \\ V(\gamma, n+1) & \text{for } A \subseteq V(\gamma, n), \ A \cap U(\gamma, n) \neq \emptyset, \ n \text{ odd} \\ V(\gamma, n) & \text{for } A \subseteq V(\gamma, n), \ A \cap U(\gamma, n) \neq \emptyset, \ n \text{ even} \\ P_{2} & \text{otherwise} \end{cases}$$

We first prove that c_1 is a closure operator.

(i) The relations $c_1(A) \subseteq P_2$ and $A \subseteq c_1(A)$ follow from the definition of c_1 .

(ii) $A' \subseteq A$ implies $c_1(A') \subseteq c_1(A)$. We distinguish some cases:

If $A = \emptyset$, then we also have $A' = \emptyset$, and $c_1(A') = c_1(A) = \emptyset$ holds.

If $A \subseteq V(\beta, n)$, $A \cap U(\beta, n) \neq \emptyset$, and *n* is even, then, we obtain by definition of $c_1, c_1(A) = V(\beta, n+1)$. Since $A' \subseteq A$, we get $A' \subseteq V(\beta, n)$. Let *r* be the maximal number such that $A' \cap U(\beta, r) \neq \emptyset$. If r = n, then we get $c_1(A') = V(\beta, n+1)$,

too, and therefore $c_1(A') = c_1(A)$. If $r \le n - 1$, then $c_1(A') \subseteq V(\beta, r')$ for some $r' \le n - 1$, which implies $c_1(A') \subset c_1(A)$.

We can analogously prove that $c_1(A') \subseteq c_1(A)$ if

 $A \subseteq V(\beta, n), A \cap U(\beta, n) \neq \emptyset$, and *n* is odd, or $A \subseteq V(\gamma, n), A \cap U(\beta, n) \neq \emptyset$, and *n* is even, or $A \subseteq V(\gamma, n), A \cap U(\beta, n) \neq \emptyset$, and *n* is odd.

If *A* is not contained in $V(\beta, n)$ and not contained in $V(\gamma, n)$ for some *n*, then $c_1(A) = P_2$. Therefore $c_1(A') \subseteq c_1(A)$ is obvious.

iii) $c_1(c_1(A)) = c_1(A)$. Again, we distinguish some cases:

If $A = \emptyset$ then $c_1(A) = \emptyset$ and $c_1(c_1(A)) = c_1(\emptyset) = \emptyset$ and $c_1(c_1(A)) = c_1(A)$ is true.

If $A \subseteq V(\beta, n)$, $A \cap U(\beta, n) \neq \emptyset$, and *n* is even, then we obtain the relation $c_1(A) = V(\beta, n+1)$. Because $c_1(A) \cap U(\beta, n+1) \neq \emptyset$ and n+1 is odd, we obtain $c_1(c_1(A)) = V(\beta, n+1)$ by definition of c_1 , which proves $c_1(c_1(A)) = c_1(A)$.

If $A \subseteq V(\beta, n)$, $A \cap U(\beta, n) \neq \emptyset$, and *n* is odd, then $c_1(A) = V(\beta, n)$. Because *n* is odd, $c_1(c_1(A)) = V(\beta, n)$, and hence $c_1(c_1(A)) = c_1(A)$.

Analogously, for $A \subseteq V(\gamma, n)$, we can prove that $c_1(c_1(A)) = c_1(A)$.

If *L* is not contained in $V(\beta, n)$ and not contained in $V(\gamma, n)$ for some *n*, then $c_1(A) = P_2$. Therefore we have $c_1(c_1(A)) = c_1(P_2) = P_2 = c_1(A)$.

We now determine the orbits of subsets of P_2 .

If $A = \emptyset$ or $A = P_2$, we get $\mathcal{O}_{c_1,d}(A) = \{A\}$. Therefore $1 \in \mathcal{O}(c_1, d)$.

If $\emptyset \subset A \subset P_2$ and A is not contained in $V(\beta, n)$ and not contained in $V(\gamma, n)$ for some n, then d(A) satisfies $\emptyset \subset d(A) \subset P_2$ and d(A) is not contained in $V(\beta, n)$ and not contained in $V(\gamma, n)$ for some n. Thus we obtain $c_1(A) = c_1(d(A)) = P_2$ and $\mathcal{O}_{c_1,d}(A) = \{A, d(A), P_2\}$. If A = d(A), e.g. for $A = D_3$, then $2 \in \mathcal{O}(c_1, d)$. If $A \neq d(A)$, e.g. for $A = \{\text{et}\}$, then $3 \in \mathcal{O}(c_1, d)$.

Let $A \subseteq V(\beta, n)$, $A \cap U(\beta, n) \neq \emptyset$, and *n* even. Starting with the closure operator c_1 , we obtain the following infinite sequences of sets:

A,

$$c_1(A) = V(\beta, n + 1),$$

 $d(c_1(A)) = V(\gamma, n + 1),$
 $c_1(d(c_1(A))) = V(\gamma, n + 2),$
 $d(c_1(d(c_1(A)))) = V(\beta, n + 2),$
 $c_1(d(c_1(d(c_1(L))))) = V(\beta, n + 3),$
 $d(c_1(d(c_1(d(c_1(L)))))) = V(\gamma, n + 3),$
....

which proves that $\mathcal{O}_{c_1,d}(A)$ is an infinite set. (For the sake of completeness we mention that the sequence starting with duality gives

A,

$$d(A),$$

 $c_1(d(A)) = V(\gamma, n),$
 $d(c_1(d(A))) = V(\beta, n)$
 $c_1(d(c_1(d(A)))) = V(\beta, n+1) = c_1(A), d(c_1(A)), ...$

which is up to first elements the same sequence which was obtained by starting with the closure operator.)

Analogously we can prove that we have infinite orbits for sets *A* where $A \subseteq V(\beta, i)$, $A \cap U(\beta, n) \neq \emptyset$, and *n* is odd or $A \subseteq V(\gamma, i)$ for some *n*.

Since we have covered all possible cases, $\mathcal{O}(c_1, d) = \{1, 2, 3, \infty\}$ follows.

Theorem 5 For any positive integer n, there is a closure operation c_2 such that $\mathcal{O}(c_2, d) = \{1, 3\} \cup \{2, 4, \dots, 2n+2\}.$

Proof Using the notation of the preceding proof, we define c_2 by

$$c_{2}(A) = \begin{cases} \emptyset & \text{for } A = \emptyset \\ V(\beta, k+1) & \text{for } A \subseteq V(\beta, k), A \cap U(\beta, k) \neq \emptyset, k \text{ even}, k < n, \\ V(\beta, k) & \text{for } A \subseteq V(\beta, k), A \cap U(\beta, k) \neq \emptyset, k \text{ even}, k \ge n, \\ V(\beta, k) & \text{for } A \subseteq V(\beta, k), A \cap U(\beta, k) \neq \emptyset, k \text{ odd}, \\ V(\gamma, k+1) & \text{for } A \subseteq V(\gamma, k), A \cap U(\gamma, k) \neq \emptyset, k \text{ odd}, k < n, \\ V(\gamma, k) & \text{for } A \subseteq V(\gamma, k), A \cap U(\gamma, k) \neq \emptyset, k \text{ odd}, k \ge n, \\ V(\gamma, k) & \text{for } A \subseteq V(\gamma, k), A \cap U(\gamma, k) \neq \emptyset, k \text{ odd}, k \ge n, \\ P_{2} & \text{otherwise} \end{cases}$$

Analogously to the proof of Theorem 4, we can show that c_2 is a closure operator.

If $A = \emptyset$ or $A = P_2$, we obtain $\mathcal{O}_{c_2,d}(A) = \{A\}$ and thus $1 \in \mathcal{O}(c_2, d)$.

If *A* is not contained in some $V(\beta, k)$ and not contained in some $V(\gamma, k), k \ge 0$, then d(A) is also not contained in some $V(\beta, k)$ and not contained in some $V(\gamma, k)$, and we obtain $\mathcal{O}_{c_2,d}(A) = \{A, d(A), P_2\}$ which gives $2, 3 \in \mathcal{O}(c_2, d)$ (as in the proof of Theorem 4).

If $A \subset V(\beta, k)$, $A \cap U(\beta, k) \neq \emptyset$, and $0 \le k < n$, then we obtain

$$\mathcal{O}_{c_{2},d}(A) = \{A, d(A)\} \cup \{V(\beta, k) \mid k \le r \le n\} \cup \{V(\gamma, k) \mid k \le r \le n\}.$$

We show this fact only for even k and even n (the proof for the other cases can be given analogously). If we start with c_2 , we get the following sets (which are obtained in succession)

$$V(\beta, k + 1), V(\gamma, k + 1), V(\gamma, k + 2), V(\beta, k + 2),$$

 $V(\beta, k + 3), V(\gamma, k + 3), \dots,$
 $V(\gamma, n - 2), V(\beta, n - 2), V(\beta, n - 1), V(\gamma, n - 1),$
 $V(\gamma, n), V(\beta, n), V(\beta, n), V(\gamma, n), V(\gamma, n), V(\beta, n), \dots,$

Deringer

and if we start with d, we get d(A), $V(\gamma, k)$, $V(\beta, k)$, $V(\beta, k+1)$ and continue as above (where we started with c_2). Therefore

$$\mathcal{N}_{c_2,d}(A) = 2(n-k+1)+2.$$

If $A = V(\beta, k)$ and $0 \le k < n$, we obtain $\mathcal{N}_{c_2,d}(A) = 2(n - k + 1)$.

If $A \subseteq V(\gamma, k)$ and k < n, $\mathcal{N}_{c_2,d}(A) \in \{2(n-k+1), 2(n-k+1)+2\}$ follows by slight modifications of the above consideration.

If $A \subset V(\beta, k)$ and $A \cap U(\beta, k) \neq \emptyset$ or $A \subset V(\gamma, k)$ and $A \cap U(\gamma, k) \neq \emptyset$ for some $k \ge n$, then we obtain $\mathcal{O}_{c_2,d}(A) = \{A, d(A), V(\beta, k), V(\gamma, k)\}$. If $A = V(\beta, k)$ or $A = V(\gamma, k)$ for some $k \ge n$, then $\mathcal{O}_{c_2,d}(A) = \{V(\beta, k), V(\gamma, k)\}$ holds.

Summarizing these facts, we obtain $\mathcal{O}(c_2, d) = \{1, 3\} \cup \{2, 4, \dots, 2n+2\}$.

Theorem 6 For any positive integer n, there is a closure operation c_3 such that $\mathcal{O}(c_3, d) = \{1, 2, 3\} \cup \{5, 7, \dots, 2n + 3\}.$

Proof We define c_3 as follows:

$$c_{3}(A) = \begin{cases} \emptyset & \text{for } A = \emptyset \\ V(\beta, k+1) & \text{for } A \subseteq V(\beta, k), A \cap U(\beta, k) \neq \emptyset, k \text{ even}, k < n, \\ V(\beta, k) & \text{for } A \subseteq V(\beta, k), A \cap U(\beta, k) \neq \emptyset, k \text{ odd}, k < n, \\ V(\gamma, k+1) & \text{for } A \subseteq V(\gamma, k), A \cap U(\gamma, k) \neq \emptyset, k \text{ odd}, k < n, \\ V(\gamma, k) & \text{for } A \subseteq V(\gamma, k), A \cap U(\gamma, k) \neq \emptyset, k \text{ even}, k < n, \\ P_{2} & \text{otherwise} \end{cases}$$

Now we follow the lines of the preceding proof; the only difference is that, for $A \subseteq V(\beta, n)$ and $A \subseteq V(\gamma, n)$ we additionally get P_2 in $\mathcal{O}_{c_3,d}(A)$.

We mention that statements analogous to the Theorems 4, 5 and 6 also hold for the involution non. We only do the following changes: Instead of β -functions we take functions from $C_2 \cap C_3$ and instead of γ -functions we take functions which are not in $C_2 \cup C_3$. Then we have the property that the negation of a function in $C_2 \cap C_3$ is not in $C_2 \cup C_3$ and vice versa.

We note that we cannot obtain arbitrary sets of natural numbers as orbits with respect to some closure operator and duality or negation. This comes from the fact that the following statement was shown in Dassow (2019): Let *c* be a closure operator on *X* and *i* an involution on *X* such that $A' \subseteq A \subseteq X$ implies $i(A') \subseteq i(A)$. If $\mathcal{K}^X(c, i) \ge 6$, then

$$\left\{2k+1 \mid 0 \le k \le \frac{\mathcal{K}^X(c,i)}{4} - 1\right\} \subset \mathcal{O}^X(c,i)$$

or

$$\left\{2k \mid 1 \le k \le \frac{\mathcal{K}^X(c,i)}{4} - 1\right\} \subset \mathcal{O}^X(c,i).$$

Consequently, because d and non satisfy the suppositions for i, certain "small" numbers have to be in $\mathcal{O}(c, d)$ and $\mathcal{O}(c, \operatorname{non})$.

5 Conclusion

In this paper we started the investigation of the Kuratowski number $\mathcal{K}(c, i)$ and the orbit $\mathcal{O}(c, i)$ for some closure operators c on the set P_2 of all Boolean functions and some involutions i on P_2 .

Especially we have determined the orbits $\mathcal{O}(S, -)$ and $\mathcal{O}(S, d)$ where *S* is the operator given by the closure defined by superposition and the involution complement or duality, respectively. For superposition and the involution negation, we have proved that $\mathcal{K}(S, \operatorname{non}) = 7$ and with respect to the orbit $\mathcal{O}(S, \operatorname{non})$, we have only left open whether or not $6 \in \mathcal{O}(S, \operatorname{non})$.

Furthermore, we have presented some closure operators c where the orbits $\mathcal{O}(c, d)$ and $\mathcal{O}(c, \operatorname{non})$ are arbitrarily large. It remains as an open problem to characterize the sets of natural numbers which can occur as orbits $\mathcal{O}(c, i)$.

Furthermore, we mention that it remains to study the Kuratowski number and the orbits if the basic set is the set P_k of all functions which map $\{0, 1, 2, ..., k-1\}^n$ into $\{0, 1, 2, ..., k-1\}$ for some n, i. e., of all functions of k-valued logic. By our results and their proofs (the functions

$$q_0(x_1, x_2) = \begin{cases} 0 & \text{for } x_1 \neq x_2 \\ x_1 + 1 & \text{mod } k & \text{for } x_1 = x_2 \end{cases}$$

and

$$q_1(x_1, x_2) = \begin{cases} 1 & \text{for } x_1 \neq x_2 \\ x_1 + 1 & \text{mod } k & \text{for } x_1 = x_2 \end{cases}$$

serve as sh and sh', respectively), we obtain

$$\mathcal{O}^{P_k}(S, \bar{}) = \{2, 4, 6\} \text{ and } \mathcal{K}(S, \bar{}) = 6.$$

Let f be a unary function of P_k such that f(f(x)) = x for all values x in $\{0, 1, 2, \dots, k-1\}$. Then we can define an f-duality d_f by

$$(d_f(g))(x_1, x_2, \dots, x_n) = f(g(f(x_1, f(x_2), \dots, f(x_n))))$$

and a "negation" $(non_f(g))(x_1, \ldots, x_n) = f(g(x_1, \ldots, x_n))$. We have no results concerning these involutions.

Funding Open Access funding enabled and organized by Projekt DEAL.

Open Access This article is licensed under a Creative Commons Attribution 4.0 International License, which permits use, sharing, adaptation, distribution and reproduction in any medium or format, as long as you give appropriate credit to the original author(s) and the source, provide a link to the Creative Commons licence, and indicate if changes were made. The images or other third party material in this article are included in the article's Creative Commons licence, unless indicated otherwise in a credit line to the material. If material is not included in the article's Creative Commons licence and your intended use is not permitted by statutory regulation or exceeds the permitted use, you will need to obtain permission directly from the copyright holder. To view a copy of this licence, visit http://creativecommons.org/licenses/by/4.0/.

References

- Brzozowski, J., Grant, E., Shallit, J.: Closures in formal languages and Kuratowski's theorem. In: Diekert, V., Nowotka, D.(eds.), Development of Languages, LNCS 5583, Springer-Verlag, pp. 125–144 (2009)
- Charlier, É., Domaratzki, M., Harju, T., Shallit, J.: Finite orbits of language operations. In: Dediu, A.-H., Martin-Vide, C., Inemaga, Sh. (eds.): Languages and Automata, Theory and Application, LNCS 6638, Springer-Verlag, pp. 204–215 (2011)
- Dassow, J.: On the orbit of closure-involution operations: the case of formal languages. Theoret. Comput. Sci. 777, 192–203 (2019)

Gardner, B.J., Jackson, M.: The Kuratowski closure-complement theorem. N. Zeal. J. Math. **38**, 9–44 (2008) Hammer, P.C.: Kuratowski's closure theorem. Nieuw Archief v. Wiskunde **7**, 74–80 (1960)

- Jablonski, S.W., Gawrilow, G.P., Kudrjawzew, W.B.: Boolesche Funktionen und Postsche Klassen. Akademie-Verlag, Berlin (1970)
- Kuratowski, K.: Sur l'opération \overline{A} de l'analysis situs. Fund. Math. **3**, 182–199 (1922)
- Post, E.L.: Introduction to a general theory of elementary propositions. Am. J. Math. 43, 163-185 (1921)
- Post, E.L.: The two-valued iterative systems of mathematical logic. Princeton University Press, Annals of Mathematics Studies (1941)

Publisher's Note Springer Nature remains neutral with regard to jurisdictional claims in published maps and institutional affiliations.